CS 70 Discrete Mathematics and Probability Theory Fall 2016 Seshia and Walrand Discussion 5B

1. **Repeated Squaring** Compute 3³⁸³ (mod 7). (Via repeated squaring!)

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Solution: Here we go...
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Divide 383 repeatedly by 2, flooring every time. We get the sequence

383, 191, 95, 47, 23, 11, 5, 2, 1.

So, to compute 3^{383} , we compute:

$$3^{1} \mod 7 \equiv 3$$

$$3^{2} \mod 7 \equiv 2$$

$$3^{5} \mod 7 \equiv (3^{2})^{2} \times 3 \equiv 2^{2} \times 3 \equiv 12 \equiv 5$$

$$3^{11} \mod 7 \equiv 5 \times 5 \times 3 \equiv 4 \times 3 \equiv 5$$

$$3^{23} \mod 7 \equiv 5 \times 5 \times 3 \equiv 5$$

$$3^{47} \mod 7 \equiv \ldots \equiv 5$$

$$3^{95} \mod 7 \equiv \ldots \equiv 5$$

$$3^{191} \mod 7 \equiv \ldots \equiv 5$$

$$3^{383} \mod 7 \equiv \ldots \equiv 5$$

2. Modular Potpourri

(a) Evaluate $4^{96} \pmod{5}$

Solution: One way: $4 \equiv -1 \pmod{5}$, and $(-1)^{96} \equiv 1$ Another: $4^2 \equiv 1 \pmod{5}$, so $4^{96} = (4^2)^{48} \equiv 1 \pmod{5}$. Mention that it is **invalid** to "apply the mod to the exponent": $4^{96} \neq 4^1 \pmod{5}$

(b) Prove or Disprove: There exists some $x \in \mathbb{Z}$ such that $x \equiv 3 \pmod{16}$ and $x \equiv 4 \pmod{6}$.

Solution: Impossible, consider both mod 2 (why is it valid to do so?)

(c) Prove or Disprove: $2x \equiv 4 \pmod{12} \iff x \equiv 2 \pmod{12}$ Solution: False, consider $x \equiv 8$.

3. Just a Little Proof

Suppose that *p* and *q* are distinct odd primes and *a* is an integer such that gcd(a, pq) = 1. Prove that $a^{(p-1)(q-1)+1} \equiv a \pmod{pq}$. **Solution:** We know that *a* is not a divsible by *p* and *a* is not divisible by *q* since gcd(a, pq) = 1.We subtract *a* from both sides to get

$$a^{(p-1)(q-1)+1} - a \equiv 0 \mod pq$$
$$a(a^{(p-1)(q-1)} - 1) \equiv 0 \mod pq$$

Since p,q are primes, we just need to show that the left hand side is divisible by both p and q. Since a is not divisible by p, we can use Fermat's Little Theorem to state that $a^{p-1} \equiv 1 \mod p$.

$$a((a^{(p-1)})^{q-1} - 1) \mod p$$

 $a(1^{q-1} - 1) \mod p$
 $0 \mod p$

Thus $a(a^{(p-1)(q-1)} - 1)$ is divisible by p. We can apply the same reasoning to show that the expression is divisible by q. Therefore we have proved our claim that $a^{(p-1)(q-1)+1} \equiv a \mod pq$.

Alternative Proof:

Because gcd(a, pq) = 1, we have that a does not divide p and a does not divide q. By Fermat's Little Theorem,

$$a^{(p-1)(q-1)+1} = (a^{(p-1)})^{(q-1)} \cdot a \equiv (1)^{q-1} \cdot a \equiv a \pmod{p}.$$

Similarly, by Fermat's Little Theorem, we have

$$a^{(p-1)(q-1)+1} = (a^{(q-1)})^{(p-1)} \cdot a \equiv (1)^{p-1} \cdot a \equiv a \pmod{q}.$$

Now, we want to use this information to conclude that $a^{(p-1)(q-1)+1} \equiv a \pmod{pq}$. We will first take a detour and show a more general result (you could write this out separately as a lemma if you want).

Consider the system of congruences

$$x \equiv a \pmod{p}$$
$$x \equiv a \pmod{q}.$$

Let's run the CRT symbolically. First off, since p and q are relatively prime, we know there exist integers g, h such that

$$g \cdot p + h \cdot q = 1.$$

We could find these via Euclid's algorithm. By the CRT, the solution to our system of congruences will be

$$x \equiv a \cdot y_1 \cdot q + a \cdot y_2 \cdot p \pmod{pq}.$$

To solve for y_1 and y_2 , we must find y_1 such that

$$x_1 \cdot p + y_1 \cdot q = 1$$

and y_2 such that

$$x_2 \cdot q + y_2 \cdot p = 1.$$

This is easy since we already know $g \cdot p + h \cdot q = 1$: the answers are $y_1 = h$ and $y_2 = g$. Finally we can plug in to the solution to get

$$x \equiv a \cdot h \cdot q + a \cdot g \cdot p \equiv a(h \cdot q + g \cdot p) \equiv a(1) \equiv a \pmod{pq}$$

Therefore by the CRT we know that the set of solutions that satisfy both $x \equiv a \pmod{p}$ and $x \equiv a \pmod{q}$ is exactly the set of solutions that satisfy $x \equiv a \pmod{pq}$.

So since $a^{(p-1)(q-1)+1} \equiv a \pmod{p}$ and $a^{(p-1)(q-1)+1} \equiv a \pmod{q}$, then by the CRT we know that $a^{(p-1)(q-1)+1}$ satisfies $a^{(p-1)(q-1)+1} \equiv a \pmod{pq}$.

4. RSA Warm-Up

Consider an RSA scheme modulus N = pq, where p and q are prime numbers larger than 3.

(a) Recall that *e* must be relatively prime to p-1 and q-1. Find a condition on *p* and *q* such that e = 3 is a valid exponent.

Solution: Both *p* and *q* must be of the form 3k + 2. p = 3k + 1 is a problem since then p - 1 has a factor of 3 in it. p = 3k is a problem because then *p* is not prime.

(b) Now suppose that p = 5, q = 17, and e = 3. What is the public key?

Solution: $N = p \cdot q = 85$ and e = 3 are displayed publically. Make sure to point out that in practice, p and q should be much larger 512-bit numbers. We are only choosing small numbers here to allow manual computation.

(c) What is the private key?

Solution: We must have $ed = 3d \equiv 1 \mod 64$, so d = 43. Reminder: we would do this by using extended gcd with x = 64 and y = 3. We get gcd(x, y) = 1 = ax + by, and a = 1, b = -21.

(d) Alice wants to send a message x = 10 to Bob. What is the encrypted message she sends using the public key?

Solution: We have $E(x) = x^3 \mod 85$. $100^3 \equiv 65 \mod 85$, so E(x) = 65.

(e) Alice receives the message y = 24 back from Bob. What equation would she use to decrypt the message?

Solution: We have $D(y) = y^{43} \mod 85$. $24^{43} \equiv 14 \mod 85$, so D(y) = 14.